

# Formal Characterisations of *Algorithm*

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# Talk Overview

- 1 Computability theory
- 2 Difficulties in accounting for algorithms
- 3 Sketch of a new approach

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- This approach has been very effective!

## Computability theory doesn't deal with algorithms

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- The implementation/algorithm distinction is completely irrelevant for computability theory
  - The distinction is also irrelevant for complexity theory
- Both these disciplines deal with the existence of effective/efficient solutions to *problems*
- Formally, we can reduce talk of algorithms in computability/complexity theory to talk of programs



Modern computer science practice marks an intuitive distinction between algorithms and implementations.

- Can we make this precise?
- What claims do we make about *algorithms*?

# Claims about algorithms

- 1 Program  $X$  implements Prim's algorithm.
- 2 MergeSort and QuickSort are different sorting algorithms.
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    - Algorithms-as-Abstracts

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- Standard philosophy of mathematics questions: out of scope



# Prim's algorithm

Prim's algorithm (Khossainov and Khossainova 2012, p. 172)

*Prim*( $G, v$ ):

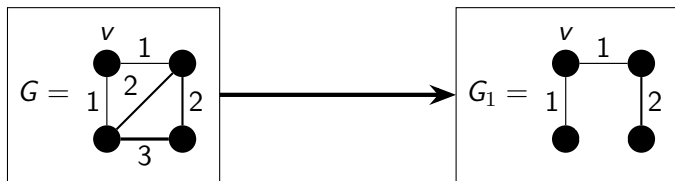
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## Testing Strong Church's Thesis

Fix a reasonable model of computation, and let  $\mathfrak{M}$  be the class of machines under that model. An *algorithm* is a machine  $M \in \mathfrak{M}$ .

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- Which  $M \in \mathfrak{M}$  is Prim's algorithm? Many issues
- Step 2a is under-determined
  - Standard models of computation force us to choose a tie-breaker

# Algorithms-as-Abstracts

- Single machines are too concrete.
- Take an equivalence relation on  $\mathfrak{M}$  - algorithms are equivalence classes
- See Dean (2007, 2016) for details on why this is dubious

# Out of options?

## Goal of project

Find a mathematical object which can adequately play the role of *algorithm*, as it's used in modern computer science practice.

- ~~Direct Algorithmic Realism~~
- ~~Strong Church's Thesis~~
- ~~Algorithms as Abstracts~~



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Two questions:

- 1 Computability theory and complexity theory don't need algorithms. Why are we assuming algorithms need them?
- 2 What do algorithms actually do?

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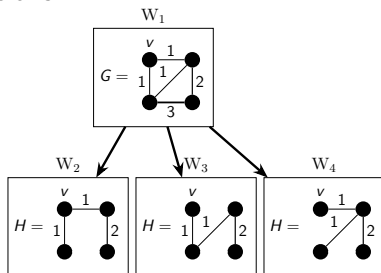
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- Let's sequester effectiveness to a requirement on *implementations*
  - This matches practice - we can describe uncomputable algorithms!

# What do algorithms do?

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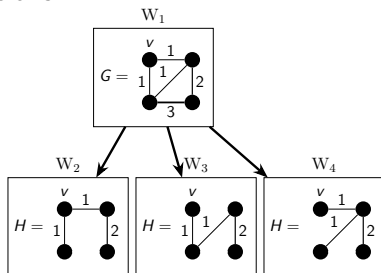
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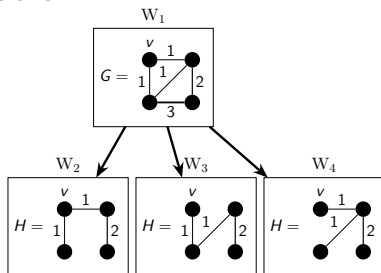
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We need an account of tasks.

- An algorithm gives you a method for achieving a task
  - Essentially, an algorithm breaks a task down into sub-tasks, or *basic operations*



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  - We can represent an algorithm “on its natural level of abstraction”
- 2 Why take algorithms' basic operations to be functions at all?
  - The tasks algorithms solve are not functions
  - If we renounce effectiveness, why stick with functions?

## Key observation

The natural basic steps of an algorithm are themselves tasks

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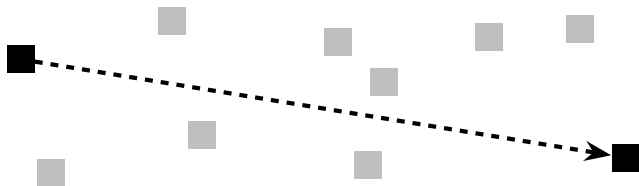
Treat tasks and basic operations as the same type of mathematical object:

- 1 We need an account of a task
- 2 We need to explain how subtasks can be combined into an algorithm
- 3 We need to explain how the resulting notion of algorithm relates to implementations

# Tasks

A task is a specification of desired behaviour.

- Formally, a task is a set of sequences of first order models
- MST is a set of pairs of first order models
- Enumerate Primes is a set of infinite sequences of first order models, each corresponding to a different enumeration of primes

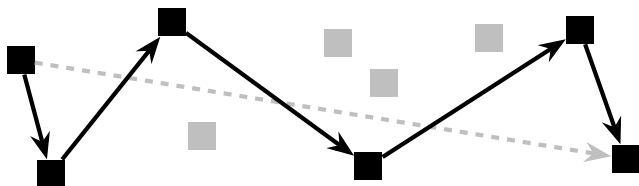




# Algorithms

An algorithm is a specification of how to solve a task, assuming the ability to solve other tasks (i.e. sub-tasks/basic operations)

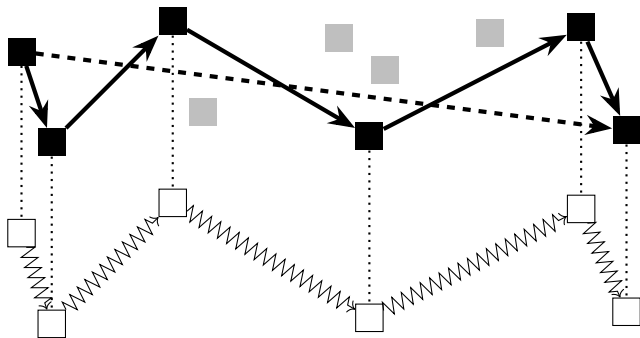
- An algorithm specifies which sub-tasks to carry out, in what order(s)
- We take an algorithm as the set of its own traces
- Formally, an algorithm is a set of sequences of first order models



# Implementations

An implementation is a simulation of the algorithm on one of the standard machine models

- Each execution trace  $s_1, s_2, \dots, s_n$  of  $M$  can be split into sub-traces corresponding to sub-tasks of the algorithm
- The sub-traces in each execution trace match the sequence of basic operations in some sequence of the algorithm



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# Summary

- Computability theory and complexity theory have no (formal) use for algorithms
- Effectiveness is a desideratum more properly suited to implementations, rather than algorithms
- If we renounce the drive for effectiveness, we can avoid problems facing many extant accounts of *algorithm*, while still maintaining compatibility with computability theory

# Thank you

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